

# Heap

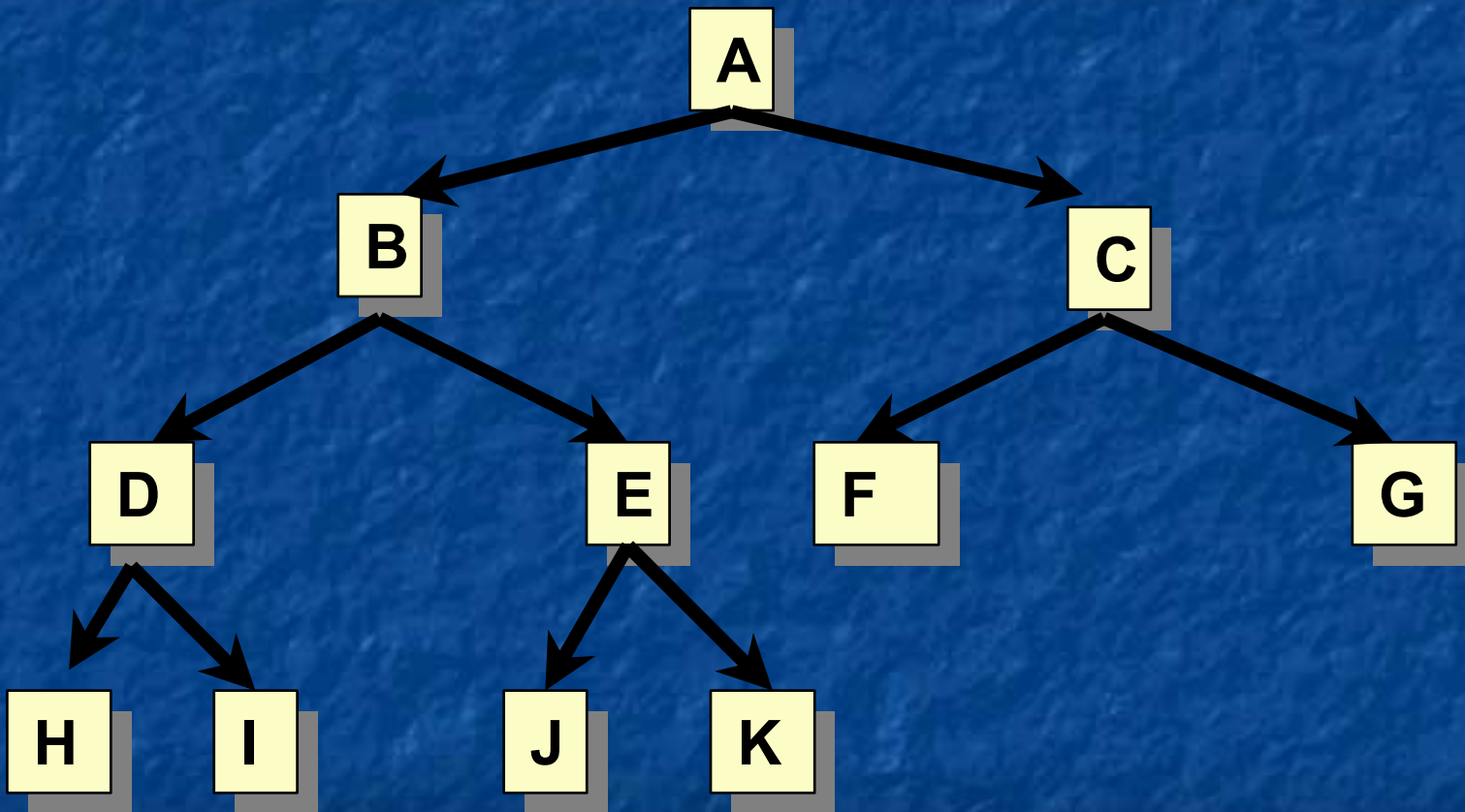
# Heap

- Definition in Data Structure
  - Heap: A special form of **complete binary tree** that key value of each node is no smaller (larger) than the key value of its children (if any).
- Max-Heap: root node has the largest key.
  - A **max tree** is a tree in which the key value in each node is **no smaller than** the key values in its children. A **max heap** is a **complete binary tree** that is also a max tree.
- Min-Heap: root node has the smallest key.
  - A **min tree** is a tree in which the key value in each node is **no larger than** the key values in its children. A **min heap** is a **complete binary tree** that is also a min tree.

# Complete Binary Tree

- A **complete binary tree** is a binary tree in which every level, *except possibly the last*, is completely filled, and all nodes are as far left as possible.

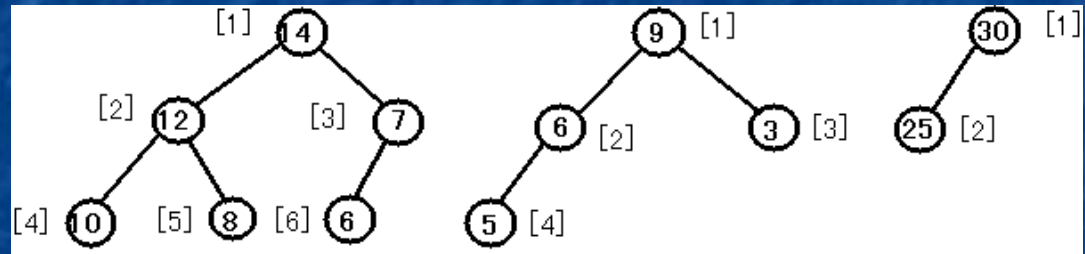
# Complete Binary Trees - Example



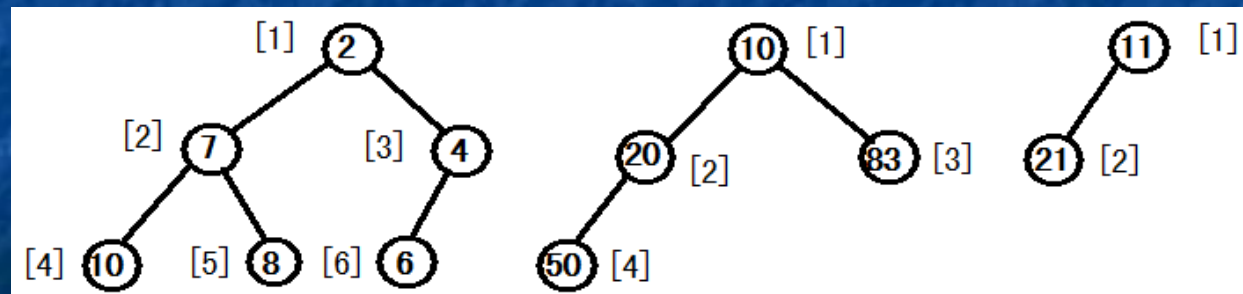
# Heap

- Example:

- Max-Heap

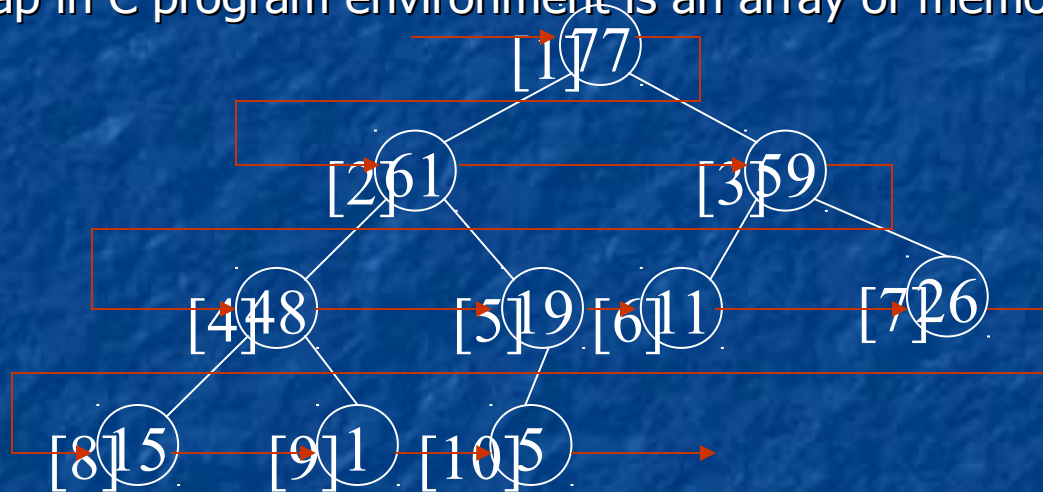


- Min-Heap



# Heap

- Notice:
  - Heap in data structure is a **complete binary tree!** (**Nice representation in Array**)
  - Heap in C program environment is an array of memory.



- Stored using array in C

index	1	2	3	4	5	6	7	8	9	10
value	77	61	59	48	19	11	26	15	1	5

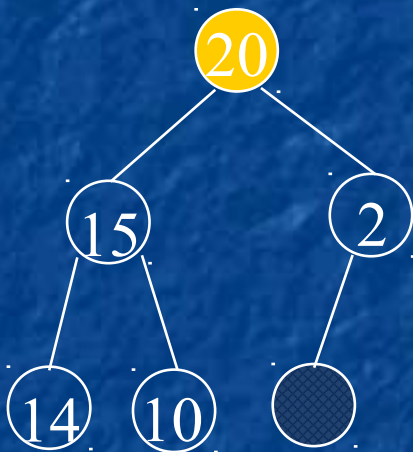
# Heap

- Operations
  - Creation of an empty heap
  - Insertion of a new element into the heap
  - Deletion of the largest(smallest) element from the heap
- Heap is complete binary tree, can be represented by array. So the complexity of inserting any node or deleting the root node from Heap is  $O(\text{height}) = O(\log_2 n)$ .

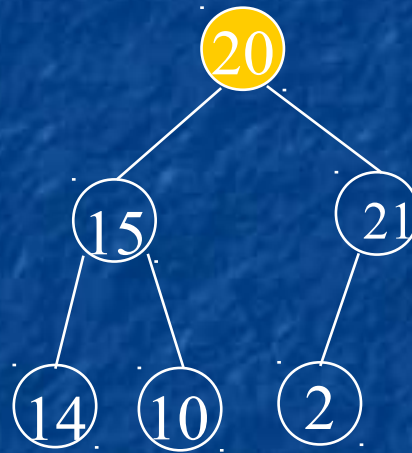
# Heap

- Given the index  $i$  of a node
- $\text{Parent}(i)$ 
  - return  $i/2$
- $\text{LeftChild}(i)$ 
  - return  $2i$
- $\text{RightChild}(i)$ 
  - Return  $2i+1$

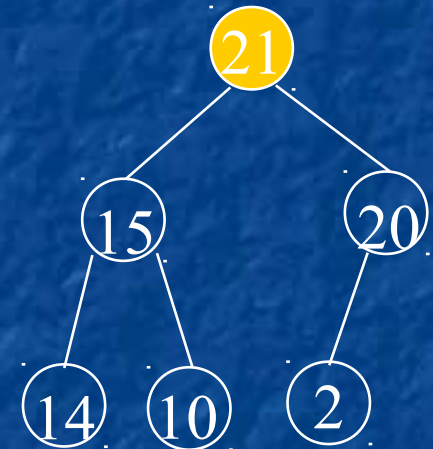
# Example of Insertion to Max Heap



initial location of new node



insert 21 into heap

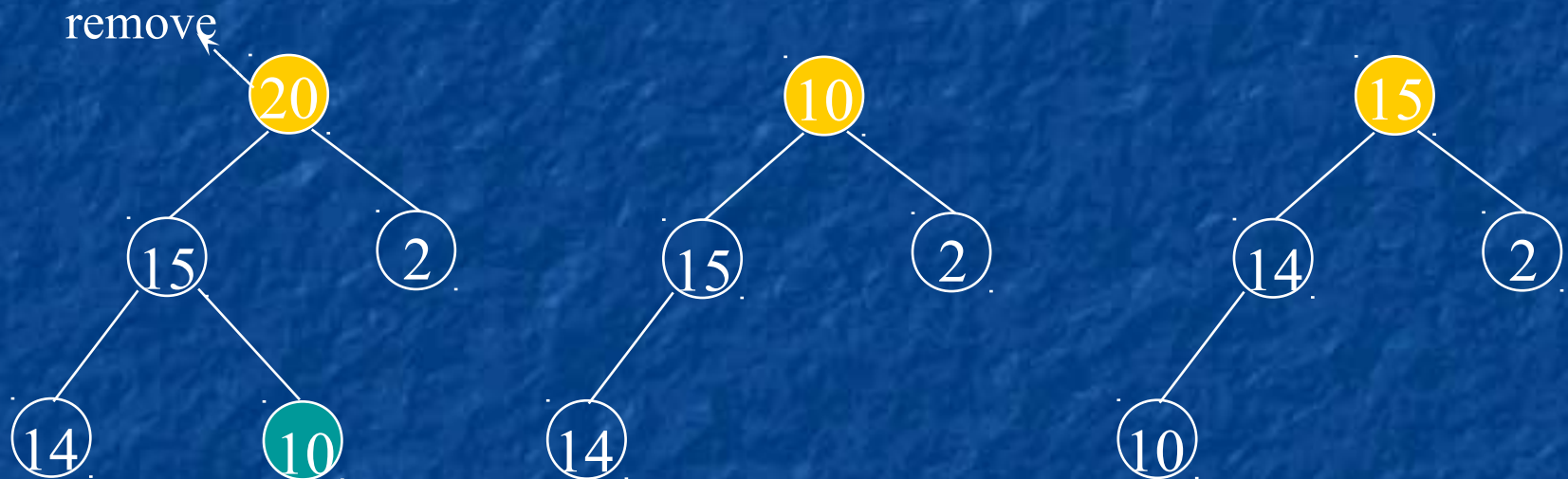


insert 21 into heap

# Insertion into a Max Heap

```
void insert_max_heap(element item, int *n)
{
    int i;
    if (HEAP_FULL(*n)) {
        fprintf(stderr, "the heap is full.\n");
        exit(1);
    }
    i = ++(*n);
    while ((i!=1)&&(item.key>heap[i/2].key)) {
        heap[i] = heap[i/2];
        i /= 2;
    }
    heap[i]= item;
}
```

# Example of Deletion from Max Heap



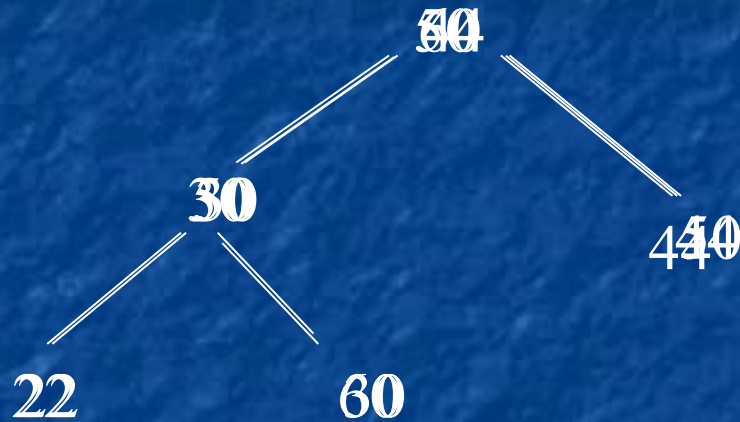
# Deletion from a Max Heap

```
element delete_max_heap(int *n)
{
    int parent, child;
    element item, temp;
    if (HEAP_EMPTY(*n)) {
        fprintf(stderr, "The heap is empty\n");
        exit(1);
    }
    /* save value of the element with the
       highest key */
    item = heap[1];
    /* use last element in heap to adjust heap */
    temp = heap[(*n)--];
    parent = 1;
    child = 2;
```

# Deletion from a Max Heap

```
while (child <= *n) {
    /* find the larger child of the current
    parent */
    if ((child < *n) &&
        (heap[child].key < heap[child+1].key))
        child++;
    if (temp.key >= heap[child].key) break;
    /* move to the next lower level */
    heap[parent] = heap[child];
    parent = child
    child *= 2;
}
heap[parent] = temp;
return item;
}
```

Suppose we want to build a heap H from the following list of numbers  
44, 30, 50, 22, 60, 55, 77, 55



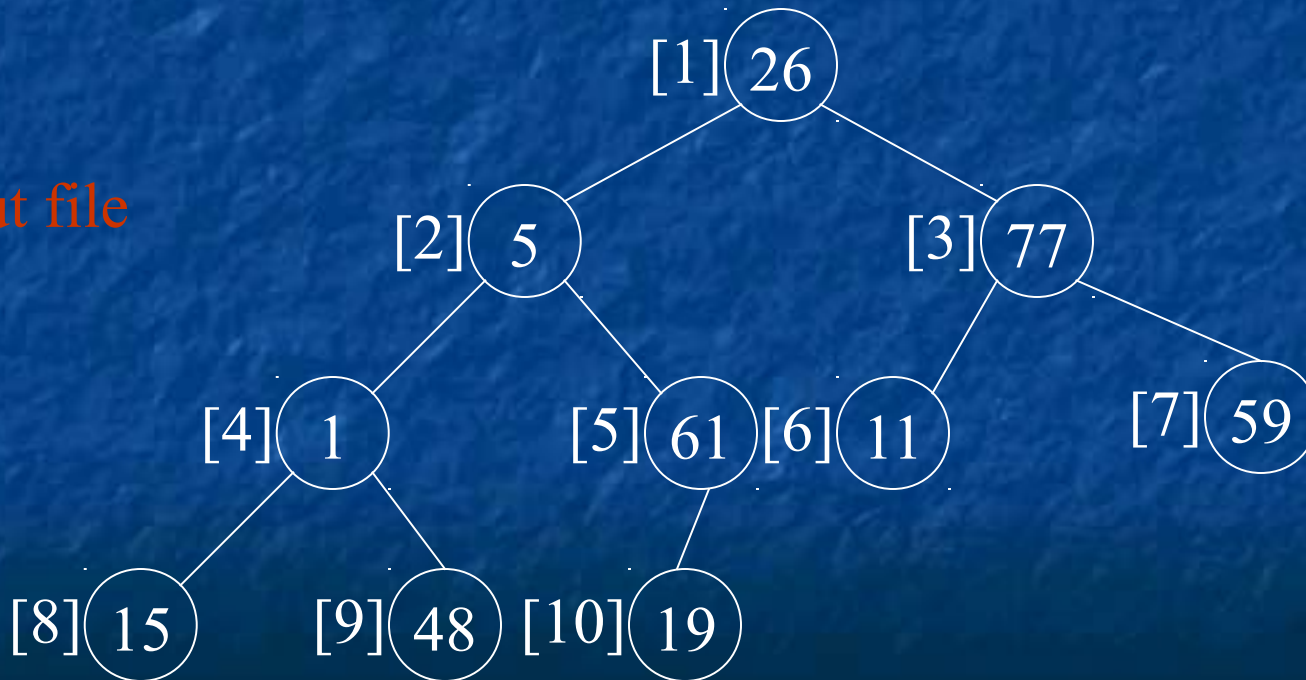


# Application On Sorting: Heap Sort

- See an illustration first
  - Array interpreted as a binary tree

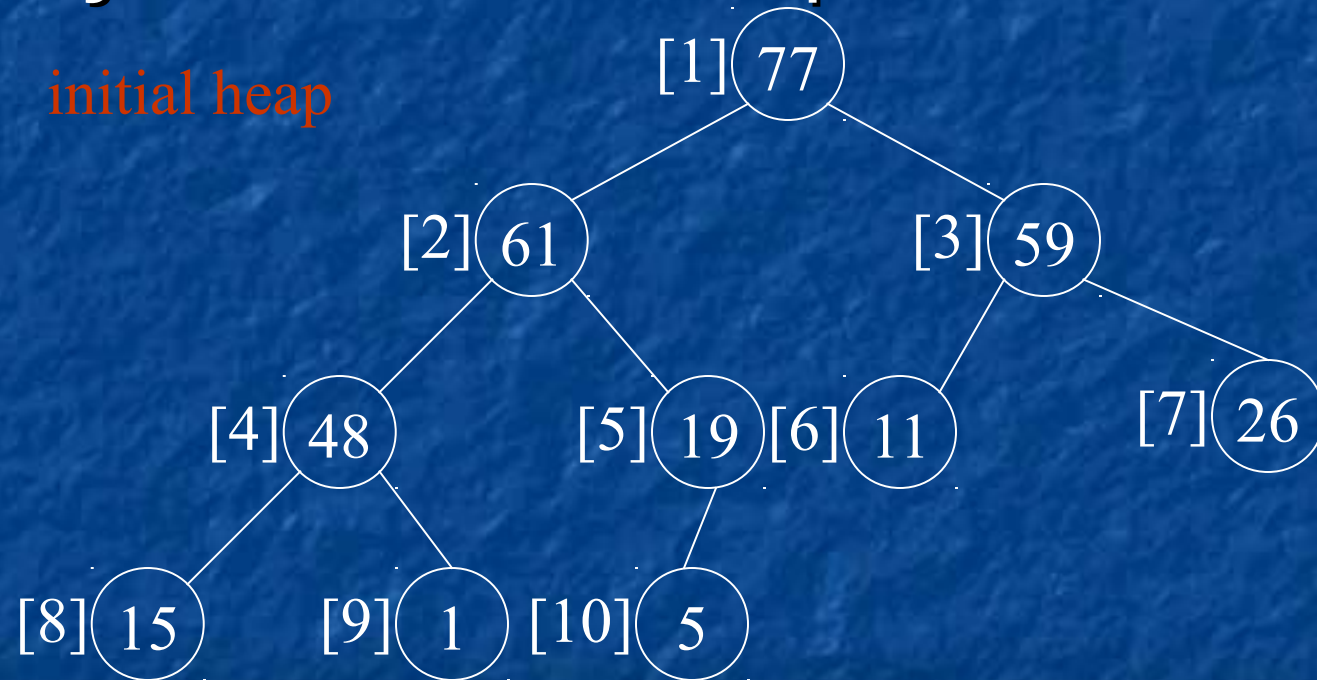
1	2	3	4	5	6	7	8	9	10
26	5	77	1	61	11	59	15	48	19

input file



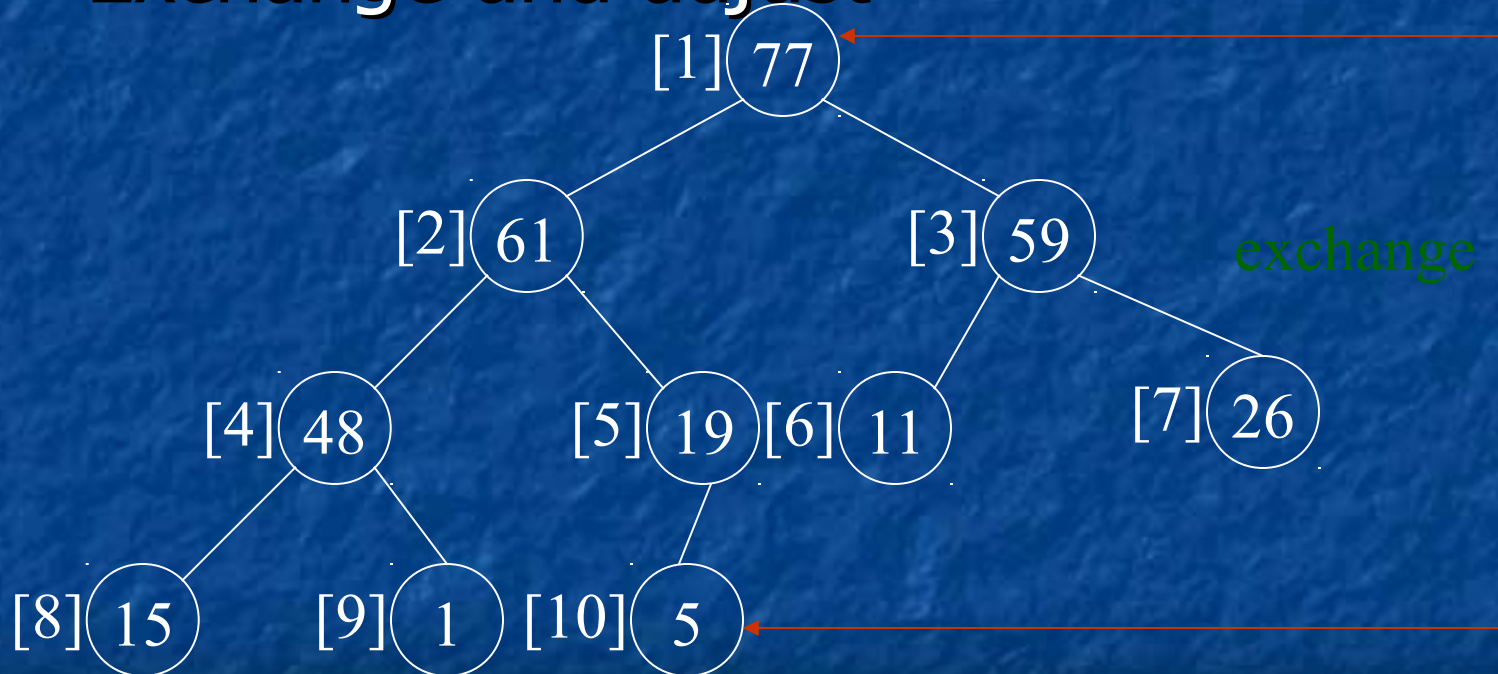
# Heap Sort

- Adjust it to a MaxHeap

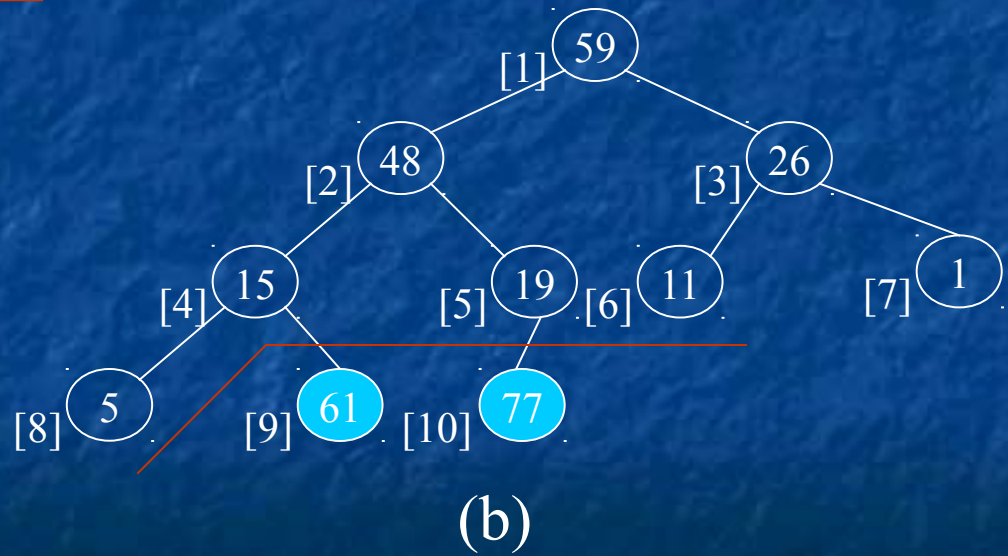
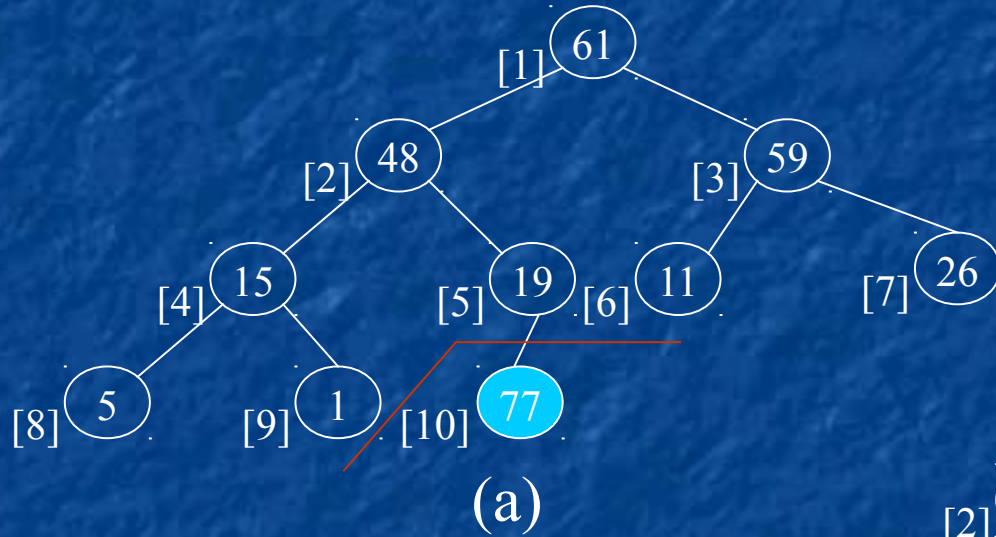


# Heap Sort

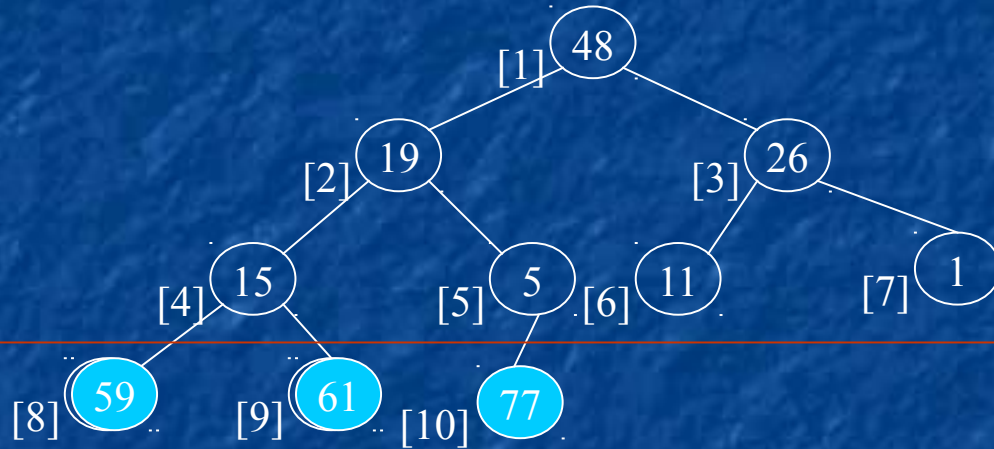
- Exchange and adjust



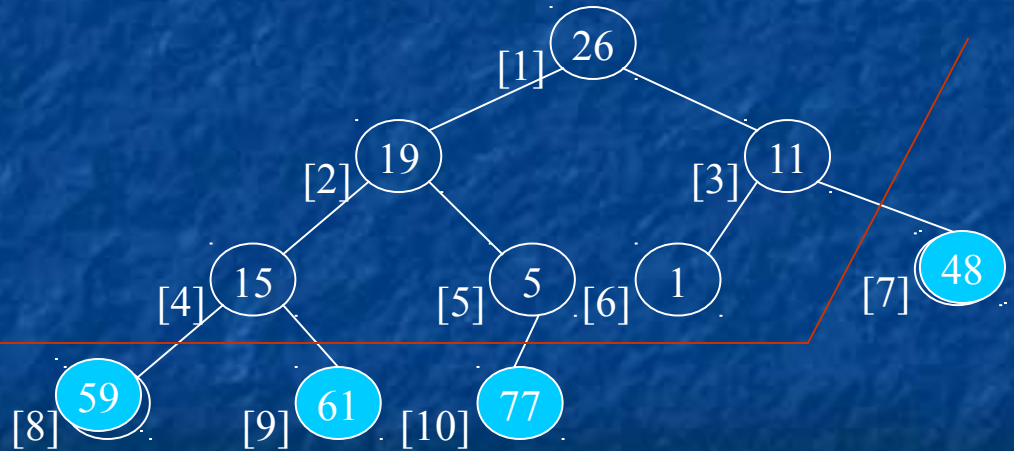
# Heap Sort



# Heap Sort

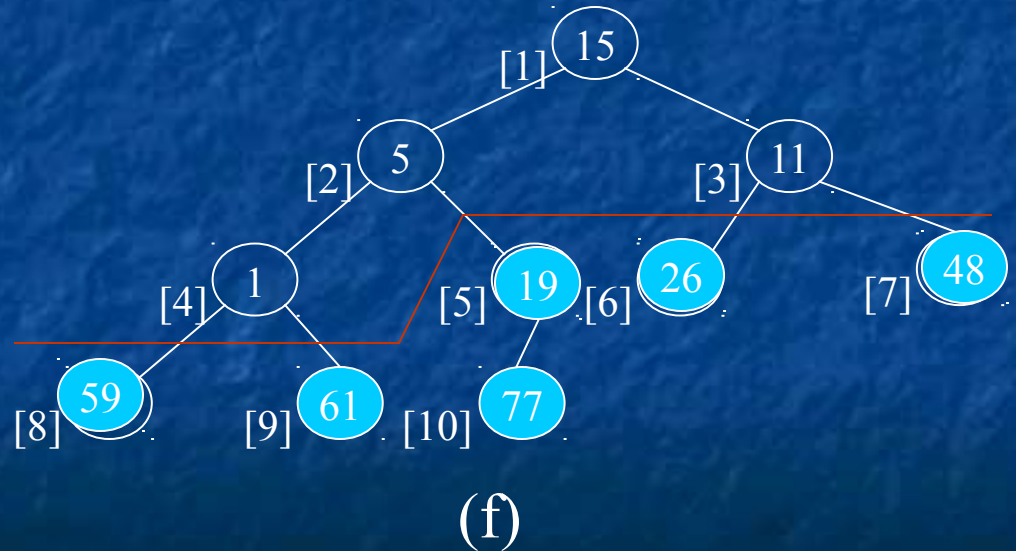
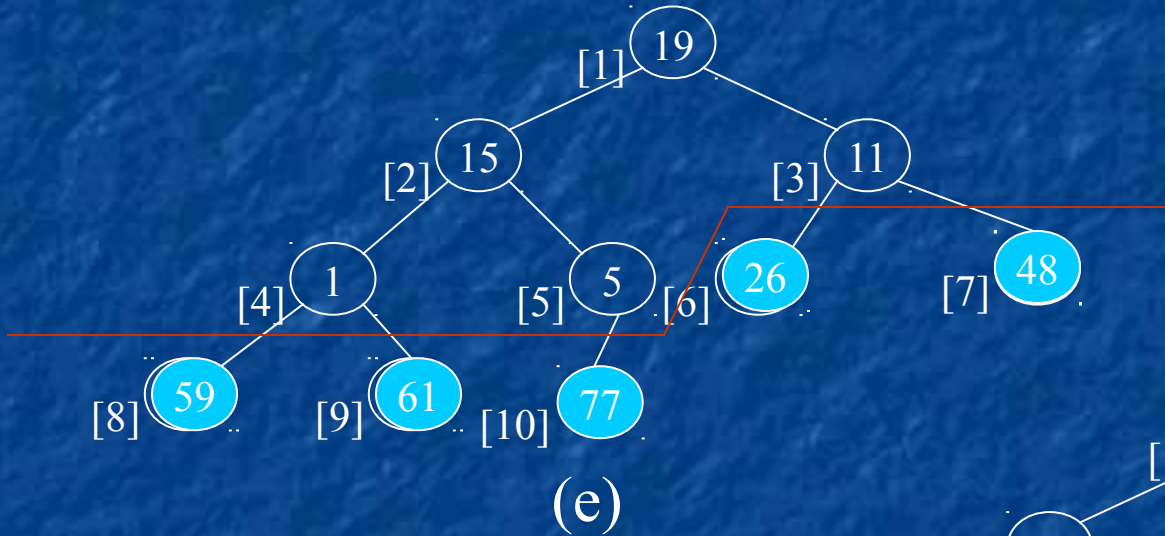


(c)

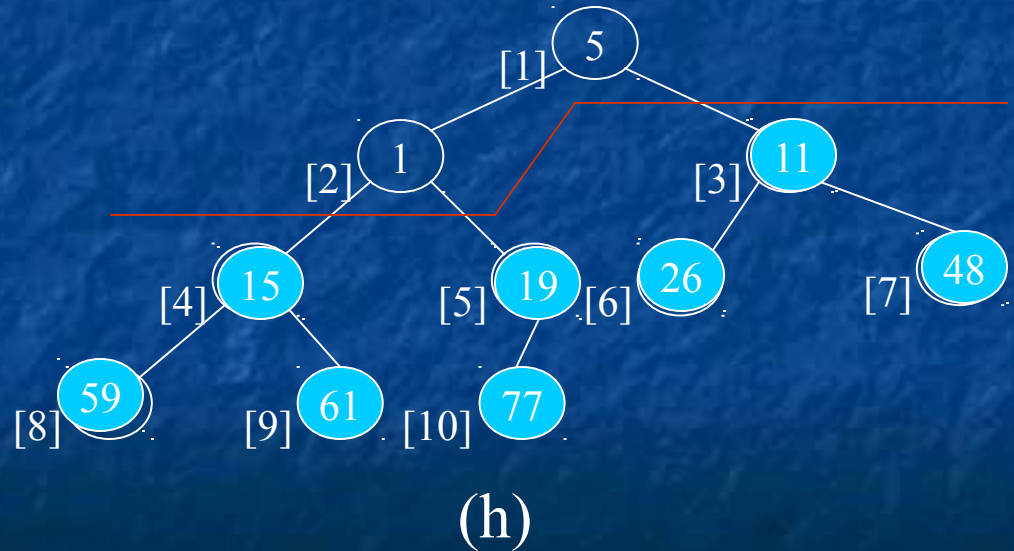
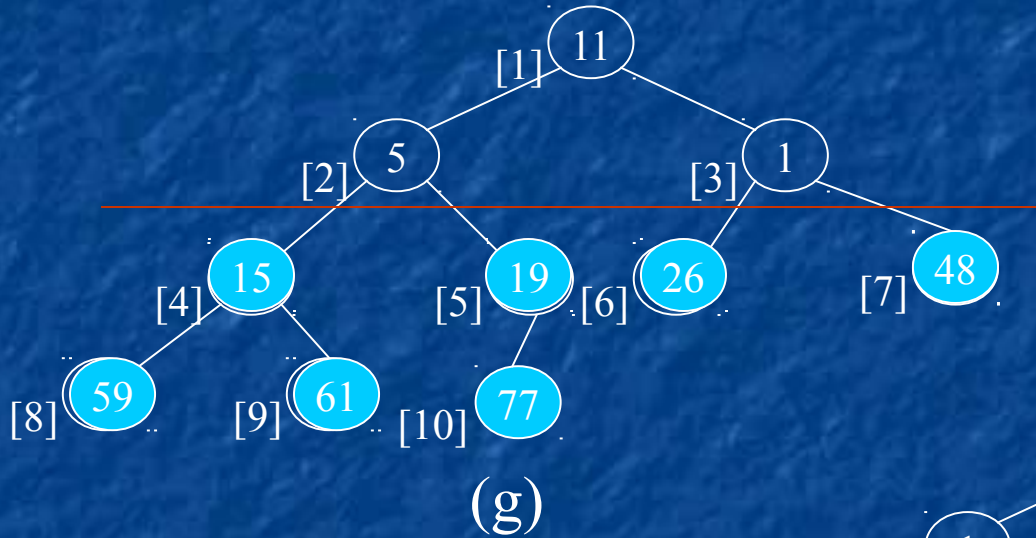


(d)

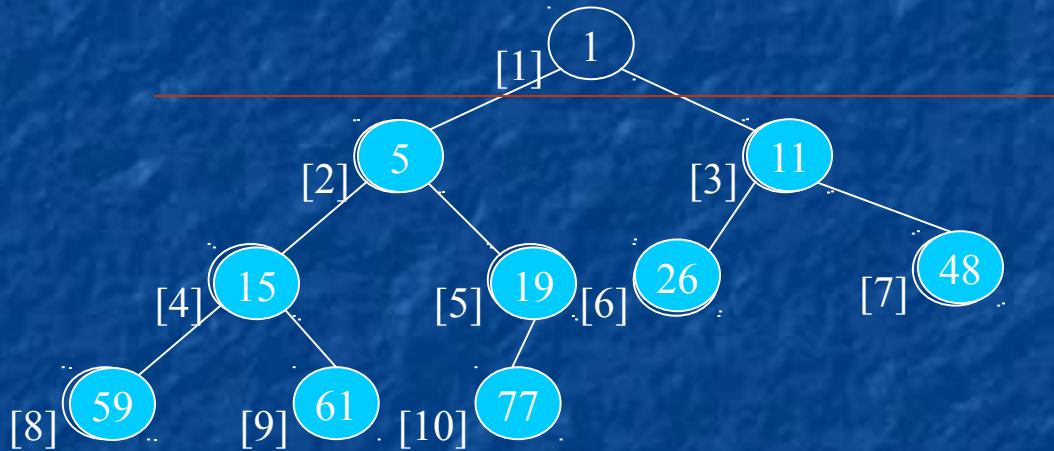
# Heap Sort



# Heap Sort



# Heap Sort



- So results (i)

77 61 59 48 26 19 15 11 5 1

# Questions

